	linnest 1	Explanation:
-1)	Every graph has an even number of odd degree vertices.	Handshaking demma: 2 deg(v) = 2.1E)
	M Webr	vev ver
	C Falsch	odd · odd = odd · even + even = even
2)	Every articulation point in a graph is incident to a bridge.	Counterexample:
	○ Wahr	Articulation polo-1
		: not a bridge
3)	There is a 5-connected graph with all vertices of degree exactly 3.	
		(Vertex-) Edge - minimum
	V Falsch	Connectivity - Connectivity - Degree
կ)	A Hamiltonian cycle in a graph visits every vertex exactly once.	Def of hamiltonian cycle.
	X Wahr	
	O Falsch	
5)	For two edges $a, b$ of a graph, the relation $a \sim b$ when $a$ and $b$ are on a common cycle is an equivalence relation.	$\begin{array}{c} \hline Def of equivalence \\ \hline \\ $
	X Wahr	relation
	O Falsch	
		Settle: En nammenhingveder Cores G = VC B websit eine Euterson
6)	Every connected graph with all vertices of even degree has an Eulerian cycle.	Euler's theorem : COULD at Inter
	X Wahr	Chas an Every vertex has an
	O Falsch	Gide Gisconnected
Ŧ)	Every 2-connected graph has a Hamiltonian cycle.	Counterex: Petersen G
	O Wahr	3-connected (thus certainly
	X Falsch	
		Goes NUT have a naminonion y
8)	If $G=(V,E)$ is a 3-connected graph and $v\in V,$ then $G[V\setminus\{v\}]$ is 2-connect	ted. Def of k-connectivity
	N Mohr	3-connected: up have to remove at
	O Falsch	I least 3 vertices to make (7)
		2- connected : " at least 2 "
3)	For every t $\geq$ 3 a complete graph K <sub>t</sub> on t vertices is 2-connected.	Def of KL: Even par of distinct vertice
	X0 Wahr	is directly joined by an edge
	O Falsch	remove 1 vertex from K,
		resulting graph is Ken

